CS 3650 Computer Systems – Summer 1 2025

Assembly

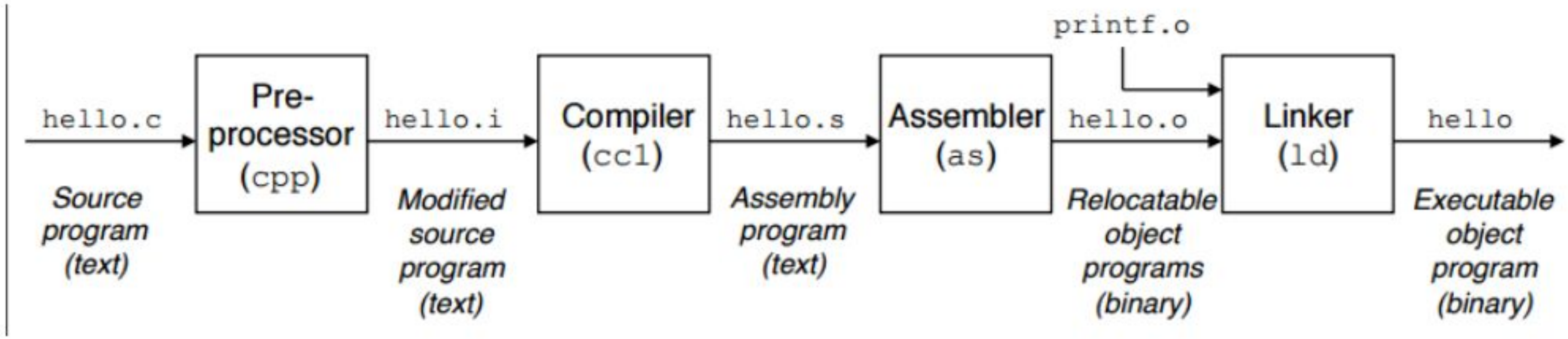
Unit 2

Northeastern University

* Acknowledgements: created based on Christo Wilson, Ferdinand Vesely, and Alden Jackson’s lecture slides for the same course. Inspired heavily by slides from Ji-Yong Shin

Recall the C toolchain pipeline

- All C programs go through this transformation of C --> Assembly --> Machine Code



So we have gone back in time in a way!

https://en.wikipedia.org/wiki/Timeline_of_programming_languages

Year	Name	Chief developer, company
1948	Curry notation system	Haskell Curry
1948	Plankalkül (concept published)	Konrad Zuse
1949	Short Code	John Mauchly and William F. Schmitt
Year	Name	Chief developer, company

1950s [edit]

Year	Name	Chief developer, company	Predecessor(s)
1950	Short Code	William F Schmidt, Albert B. Tonik,[3] J.R. Logan	Brief Code
1950	Birkbeck Assembler	Kathleen Booth	ARC
1951	Superplan	Heinz Rutishauser	Plankalkül
1951	ALGAE	Edward A Voorhees and Karl Balke	none (unique language)
1951	Intermediate Programming Language	Arthur Burks	Short Code
1951	Regional Assembly Language	Maurice Wilkes	EDSAC
1951	Boehm unnamed coding system	Corrado Böhm	CPC Coding scheme
1951	Klammerausdrücke	Konrad Zuse	Plankalkül
1951	OMNIBAC Symbolic Assembler	Charles Katz	Short Code
1951	Stanislaus (Notation)	Fritz Bauer	none (unique language)
1951	Whirlwind assembler	Charles Adams and Jack Gilmore at MIT Project Whirlwind	EDSAC
1951	Rochester assembler	Nat Rochester	EDSAC

So we have gone back in time!

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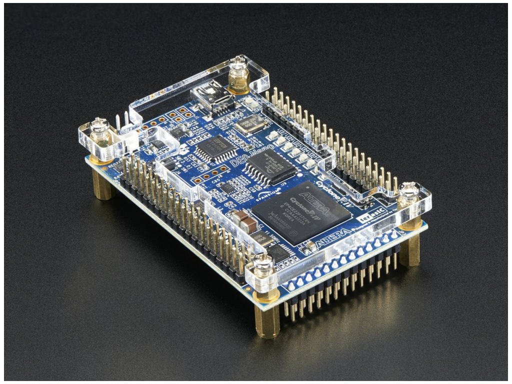
Look at all of these assembly languages over 60 years old!

This was the family of languages folks programmed in.

Modern Day Assembly is of course still in use

- Still used in games
  (console games specifically)
  - In hot loops where code must run fast
- Still used on embedded systems
- Useful for debugging any compiled language
- Useful for even non-compiled or Just-In-Time Compiled languages
  - Python has its own bytecode
  - Java's bytecode (which is eventually compiled) is assembly-like
- Being used on the web
  - webassembly
- Still relevant after 60+ years!



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Aside: Java(left) and Python(right) bytecode examples

```
 0 aload_0
 1 new #3 <acceptanceTests/treeset_personOK/Main$A>
 4 dup
 5 new #8 <java/lang/Object>
 8 dup
 9 invokespecial #10 <java/lang/Object.<init>>
12 new #12 <java/lang/Integer>
15 dup
16 iconst_2
17 invokespecial #14 <java/lang/Integer.<init>>
20 invokespecial #17 <acceptanceTests/treeset_personOK/Main$A.<init>>
23 new #12 <java/lang/Integer>
26 dup
27 iconst_1
28 invokespecial #14 <java/lang/Integer.<init>>
31 invokespecial #17 <acceptanceTests/treeset_personOK/Main$A.<init>>
34 getstatic #20 <java/lang/System.out>
37 new #3 <acceptanceTests/treeset_personOK/Main$A>
40 dup
41 new #8 <java/lang/Object>
44 dup
45 invokespecial #10 <java/lang/Object.<init>>
48 new #12 <java/lang/Integer>
51 dup
52 iconst_2
53 invokespecial #14 <java/lang/Integer.<init>>
56 invokespecial #17 <acceptanceTests/treeset_personOK/Main$A.<init>>
59 invokevirtual #26 <java/io/PrintStream.println>
62 return
```

```
def f(n, accum):
  if n <= 1:
    return accum
  else:
    return f(n-1, accum*n)
```

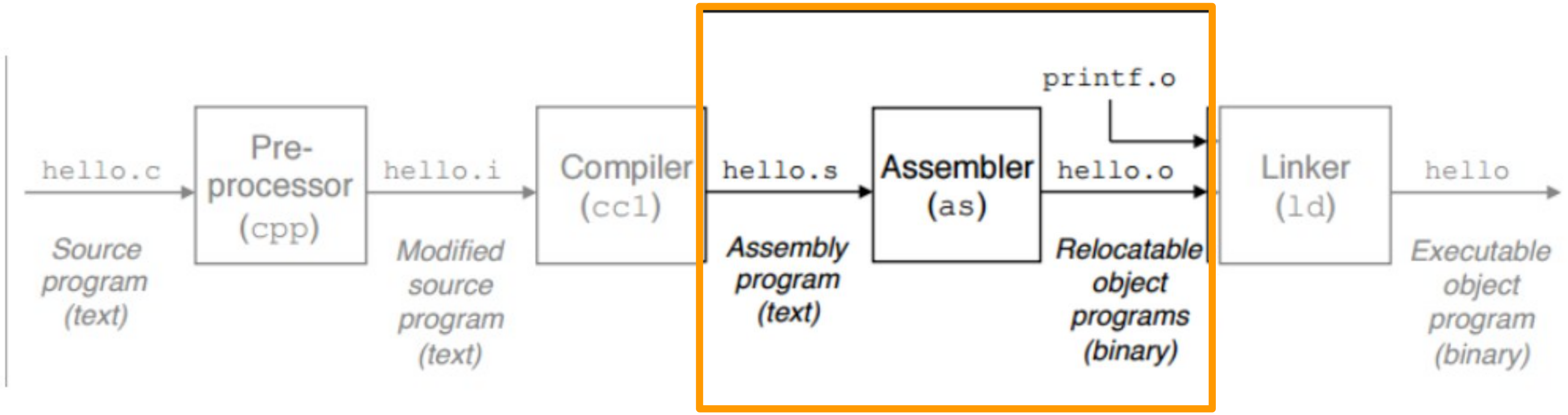
```
>>> import dis
>>> dis.dis(f)
  2           0 LOAD_FAST                0 (n)
              3 LOAD_CONST               1 (1)
              6 COMPARE_OP               1 (<=)
              9 POP_JUMP_IF_FALSE       16

  3          12 LOAD_FAST                1 (accum)
             15 RETURN_VALUE

  5     >>   16 LOAD_GLOBAL              0 (f)
             19 LOAD_FAST                0 (n)
             22 LOAD_CONST               1 (1)
             25 BINARY_SUBTRACT
             26 LOAD_FAST                1 (accum)
             29 LOAD_FAST                0 (n)
             32 BINARY_MULTIPLY
             33 CALL_FUNCTION            2
             36 RETURN_VALUE
             37 LOAD_CONST               0 (None)
             40 RETURN_VALUE
```

Assembly is important in our toolchain

- Even if the step is often hidden from us!

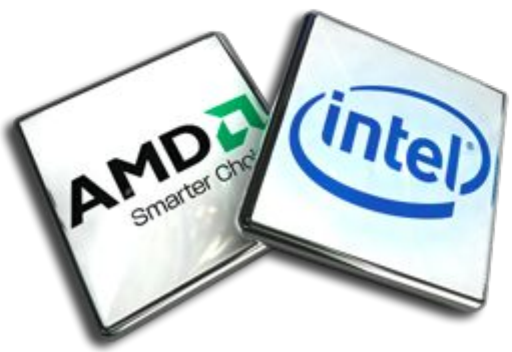


Intel and x86 Instruction set

- In order to program these chips, there is a specific instruction set we will use
- Popularized by Intel
- Other companies have contributed.
    - AMD has been the main competitor
- (AMD was first to really nail 64 bit architecture around 2001)
- Intel followed up a few years later (2004)
- Intel remains the dominant architecture
- x86 is a CISC architecture
    - (CISC pronounced /ˈsɪsk/)

CISC versus RISC

- Complex Instruction Set Computer (CISC)
  - Instructions do more per operation
  - Architecture understands a series of operations
- Performance can be nearly as fast or equal to RISC



- Reduced Instruction Set Computer (RISC)
  - Instructions are very small
  - Performance is extremely fast
  - Generally a simpler architecture

Introduction to Assembly

How are programs created?

- Compile a program to an executable
  - gcc main.c -o program

- Compile a program to assembly
  - gcc main.c -S -o main.s

- Compile a program to an object file (.o file)
  - gcc -c main.c

- Linker (A program called ld) then takes all of your object files and makes a binary executable.

Focus on this step today

- ~~Compile a program to an executable~~
  - ~~gcc main.c -o program~~

- Compile a program to assembly
  - gcc main.c -S -o main.s

- ~~Compile a program to an object file (.o file)~~
  - ~~gcc -c main.c~~

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Layers of Abstraction

- As a C programmer you worry about C code
  - You work with variables, do some memory management using malloc and free, etc.
- As an assembly programmer, you worry about assembly
  - You also maintain the registers, condition codes, and memory
- As a hardware engineer (programmer)
  - You worry about cache levels, layout, clocks, etc.

Assembly Abstraction layer

- With Assembly, we lose some of the information we have in C

- In higher-order languages we have many different data types which help protect us from errors.
  - For example: int, long, boolean, char, string, float, double, complex, …
  - In C there are custom data types (structs for example)
  - Type systems help us avoid inconsistencies in how we pass data around.

- In Assembly we lose unsigned/signed information as well!
  - However, we do have two data types
  - Types for integers (1,2,4,8 bytes) and floats (4,8, or 10 bytes) [byte = 8 bits]

Sizes of data types (C to assembly)

C Declaration	Intel Data Type	Assembly-code suffix	Size (bytes)
char	Byte	b	1
short	Word	w	2
int	Double word	l	4
long	Quad word	q	8
char *	Quad word	q	8
float	Single precision	s	4
double	Double Precision	l	8

Sizes of data types (C to assembly)

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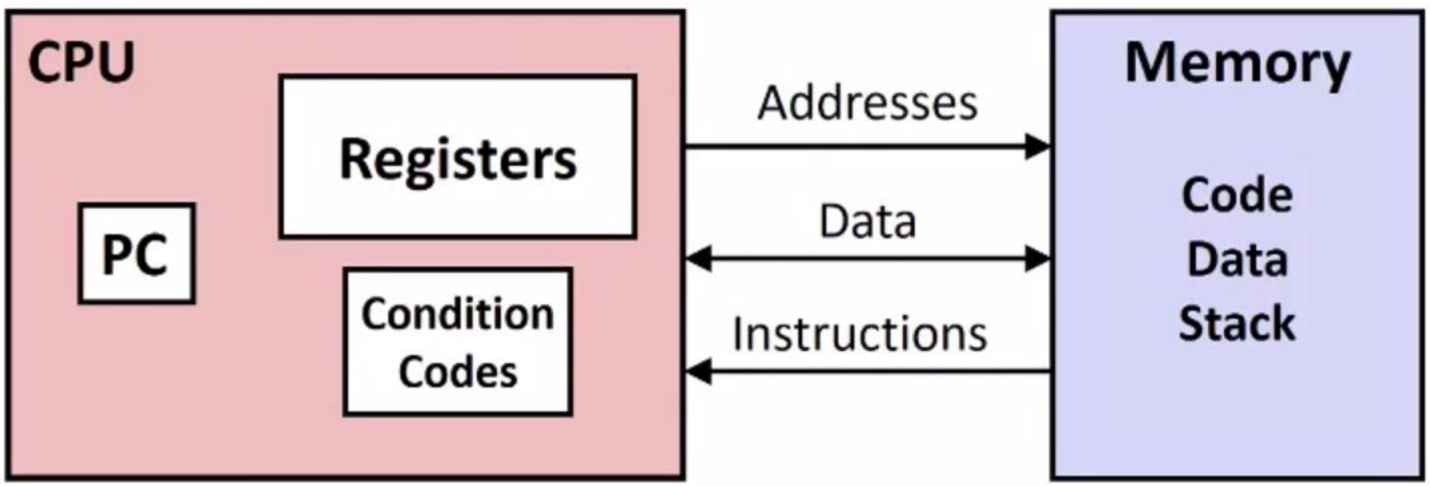
For us, one word of data is 64 bits
[8 bytes] but may vary on other hardware

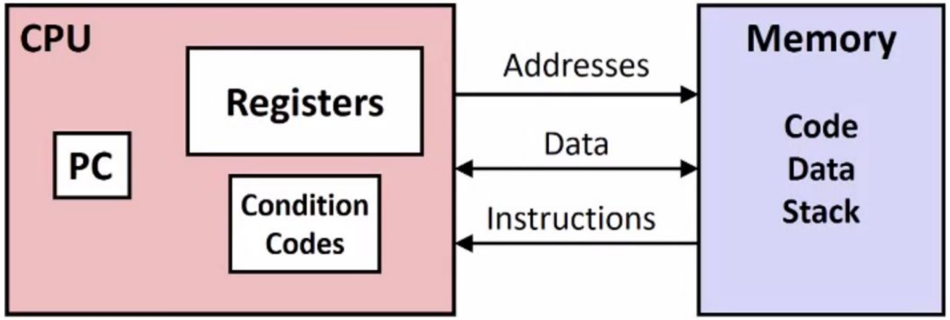
View as an assembly programmer

- Register - where we store data (heavily used data)
- PC - gives us address of next instruction
- Condition codes - some status information
- Memory – where the program (code) resides and data is stored



Assembly Operations (i.e. Our instruction set)

- Things we can do with assembly (and this is about it!)
  - Transfer data between memory and register
    - Load data from memory to register
    - Store register data back into memory
  - Perform arithmetic/logical operations on registers and memory
  - Transfer Control
    - Jumps
    - Branches (conditional statements)



x86-64 Registers

- Focus on the 64-bit column.
- These are 16 general purpose registers for storing bytes
  - (Note sometimes we do not always have access to all 16 registers)
- Registers are similar to variables where we store values

Not modified for 8-bit operands

Not modified for 16-bit operands

Zero-extended for 32-bit operands

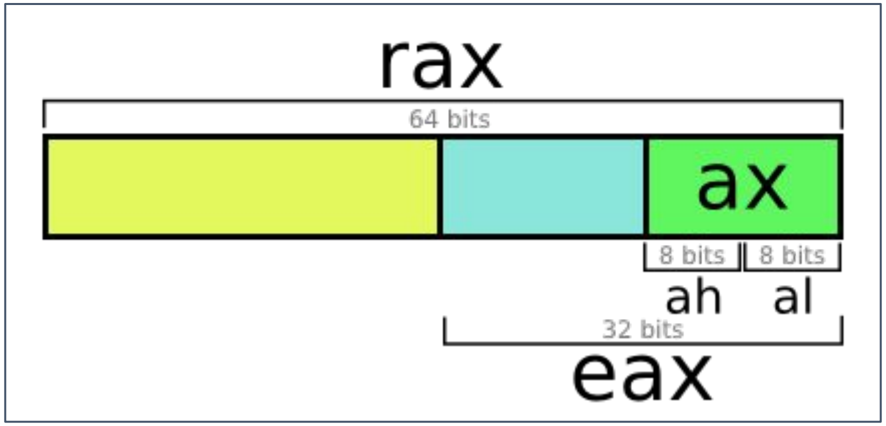
Register encoding		Low 8-bit	16-bit	32-bit	64-bit
0	AH†	AL	AX	EAX	RAX
3	BH†	BL	BX	EBX	RBX
1	CH†	CL	CX	ECX	RCX
2	DH†	DL	DX	EDX	RDX
6		SIL‡	SI	ESI	RSI
7		DIL‡	DI	EDI	RDI
5		BPL‡	BP	EBP	RBP
4		SPL‡	SP	ESP	RSP
8		R8B	R8W	R8D	R8
9		R9B	R9W	R9D	R9
10		R10B	R10W	R10D	R10
11		R11B	R11W	R11D	R11
12		R12B	R12W	R12D	R12
13		R13B	R13W	R13D	R13
14		R14B	R14W	R14D	R14
15		R15B	R15W	R15D	R15

63 32 31 16 15 8 7 0

† Not legal with REX prefix ‡ Requires REX prefix

x86-64 Register (zooming in)

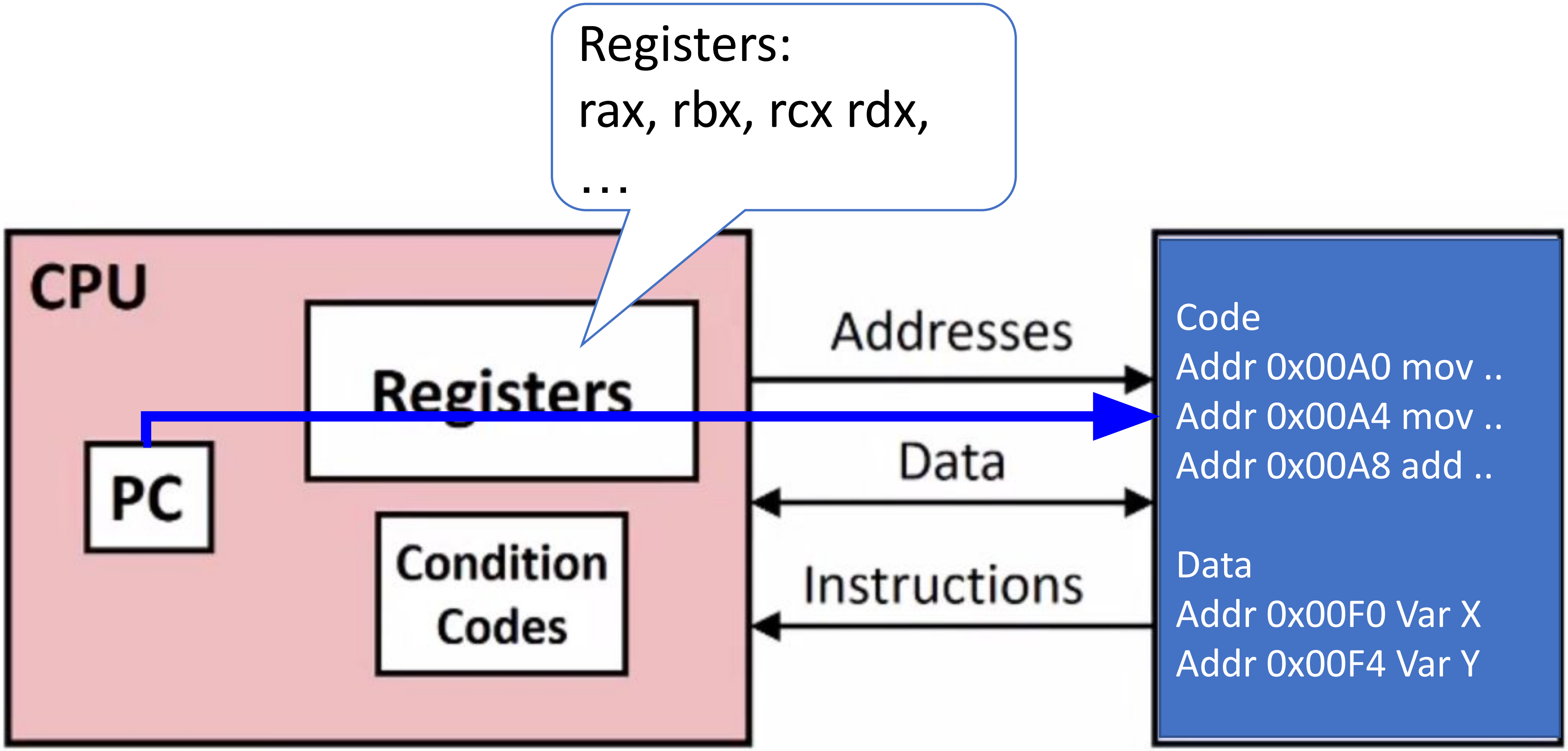
- Note register eax addresses the lower 32 bits of rax
- Note register ax addresses the lower 16 bits of eax
- Note register ah addresses the high 8 bits of ax
- Note register al (lowercase L) addresses the low 8 bits of ax



Some registers are reserved for special use (More to come)

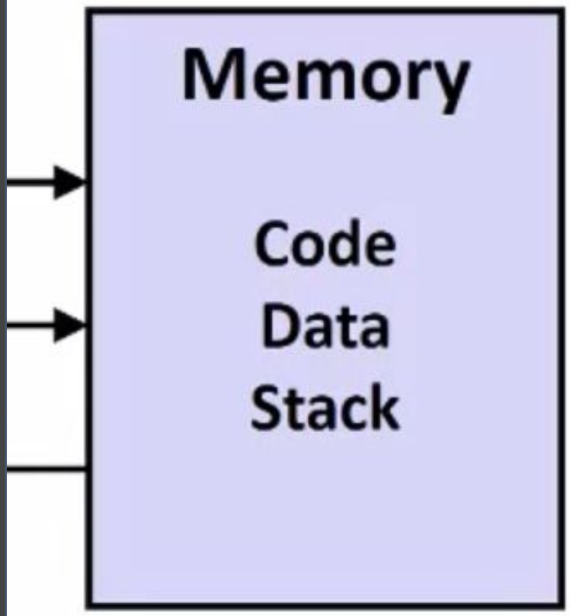
- This can be dependent on the instruction being used
  - %rsp - keeps track of where the stack pointer is
  - (We will do an example with the stack and what this means soon)

Program Counter and Memory Addresses



Memory Addresses

- Note that we are looking at virtual addresses in our assembly when we see addresses.
- This makes us think of the program as a large byte array.
  - The operating system takes care of managing this for us with virtual memory.
  - This is one of the key jobs of the operating system



A First Assembly Instruction

Moving data around | mov instruction

- (Remember moving data is all machines do!)
- movq - moves a quad word (8 bytes) of data
- movd - move a double word (4 bytes) of data

```
movq Source, Dest
```

Order matters
"source to destination"
"left to right"

Moving data around | mov instruction

- (Remember moving data is all machines do!)
- movq - moves a quad word (8 bytes) of data
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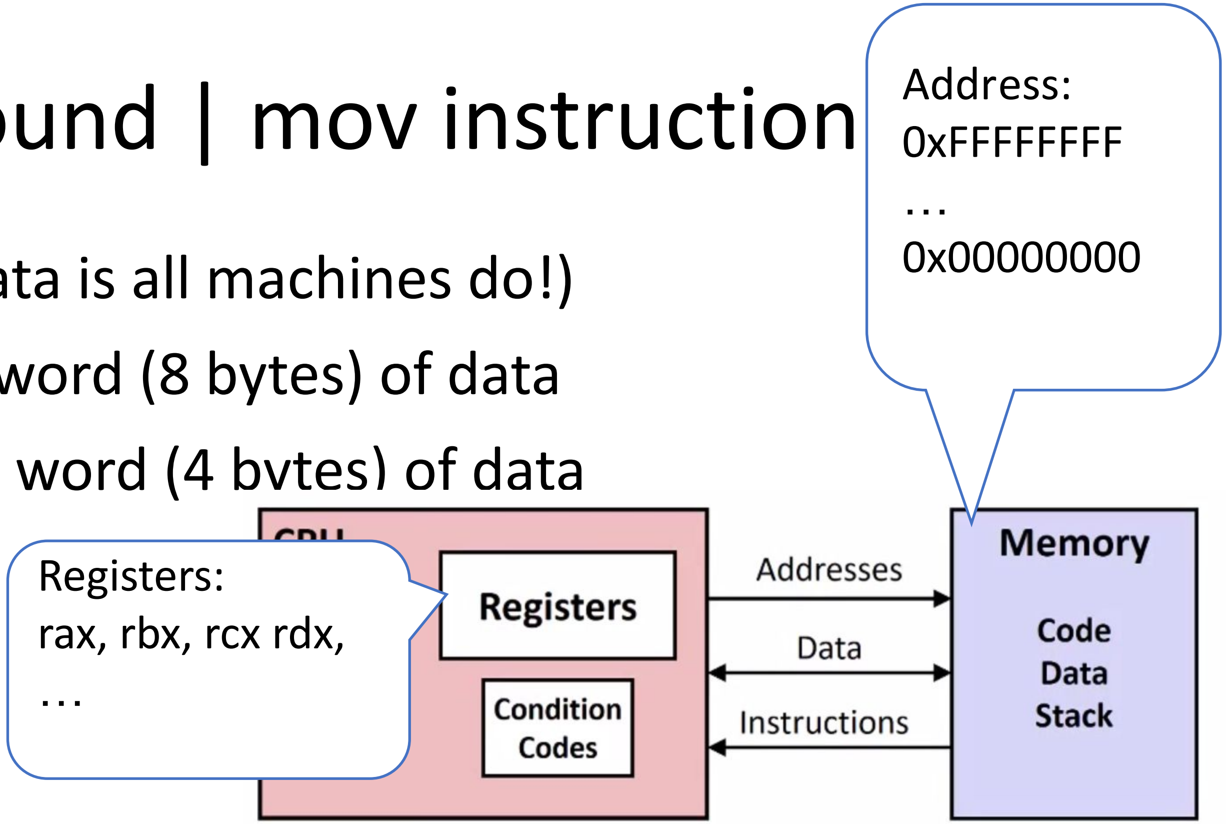
`movq Source, Dest`



- Source or Dest Operands can have different addressing modes
  - Immediate - some memory address **$**0x333 or **$**-900
  - Memory - **(**%rax**)** dereferences gets the value in the register and use it as address
  - Register - Just **%**rax

Full List of Memory Addressing Modes

Mode	Example
Global Symbol	MOVQ x, %rax
Immediate	MOVQ $56, %rax
Register	MOVQ %rbx, %rax
Indirect	MOVQ (%rsp), %rax
Base-Relative	MOVQ -8(%rbp), %rax
Offset-Scaled-Base-Relative	MOVQ -16(%rbx, %rcx, 8), %rax (*base, index, scale*)

Copy data from addr pointed by rbp minus 8 to rax

(rbx + rcx * 8) - 16

C equivalent of movq instructions | movq src, dest

```
movq $0x4, %rax
movq $-150, (%rax)

movq %rax, %rdx
movq %rax, (%rdx)

movq (%rax), %rdx
```

What does each movq do?

C equivalent of movq instructions | movq src, dest

`movq $0x4, %rax`	**`%rax`** `= 0x4;` (Moving in literal value into register)
`movq $-150, (%rax)`	use value of rax as memory location and set that location to -150 (`*p = -150`)
`movq %rax, %rdx`	`%rdx = %rax` (copy src into dest)
`movq %rax, (%rdx)`	use value of rdx as memory location and set that location to value stored in rax (`*p = %rax`)
`movq (%rax), %rdx`	Set value of rdx to value of rax as memory location (`%rdx = *p`)

Some registers are reserved for special use (More to come)

- This can be dependent on the instruction being used
- %rsp - keeps track of where the stack is for example
- %rdi - the first program argument in a function
- %rsi - the second argument in a function
- %rdx - the third argument of a function
- %rax – return value of a function

These conventions are especially useful for functions known as system calls.

1	write	sys_write	fs/read_write.c
%rdi **unsigned int** fd	%rsi **const char __user** * buf	%rdx **size_t** count	

https://filippo.io/linux-syscall-table/

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- **%rip - the Program Counter**

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- %rax – return value of a function
- %rip - the Program Counter
- **%r8-%r15 - These eight registers are general purpose registers**

A little example

What does this function do?
(take a few moments to think)

- void mystery(<type> a, <type> b) {

    ????

  }

- mystery:

```
movq (%rdi), %rax
movq (%rsi), %rdx
movq %rdx, (%rdi)
movq %rax, (%rsi)
ret
```

Cheat Sheet

(Note: This can be dependent on the instruction being used)
%rsp - keeps track of where the stack is for example
%rdi - the first program argument in a function
%rsi - The second argument in a function
%rdx - the third argument of a function
%rip - the Program Counter
%r8-%r15 - These ones are actually the general purpose registers

swap of long

- void mystery(long *a, long *b) {

```
    long t0 = *a;
    long t1 = *b;
    *a = t1;
    *b = t0;
}
```

- mystery:

```
    movq (%rdi), %rax
    movq (%rsi), %rdx
    movq %rdx, (%rdi)
    movq %rax, (%rsi)
    ret
```

Cheat Sheet

(Note: This can be dependent on the instruction being used)
%rsp - keeps track of where the stack is for example
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More assembly instructions

```
addq  Src, Dest      Dest=Dest+Src
subq  Src, Dest      Dest=Dest-Src
imulq Src, Dest      Dest=Dest*Src
salq  Src, Dest      Dest=Dest << Src
sarq  Src, Dest      Dest=Dest >> Src
shlq  Src, Dest      Dest=Dest << Src
shrq  Src, Dest      Dest=Dest >> Src
xorq  Src, Dest      Dest=Dest ^ Src
andq  Src, Dest      Dest=Dest & Src
orq   Src, Dest      Dest=Dest | Src
```

- Note on order:
  We use AT&T syntax: op Src, Dest
  Intel syntax: op Dest, Src

	Value 1	Value 2
x	0110 0011	1001 0101
x>>4 (arithmetic)	**0000** 0110	**1111** 1001
x>>4 (logical)	**0000** 0110	**0000** 1001

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Exercise

- If I have the expression

  c = b*(b+a)

- How should I write this is ASM?

Cheat Sheet

addq	Src, Dest	Dest=Dest+Src
subq	Src, Dest	Dest=Dest-Src
imulq	Src, Dest	Dest=Dest*Src
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shrq	Src, Dest	Dest=Dest >> Src
xorq	Src, Dest	Dest=Dest ^ Src
andq	Src, Dest	Dest=Dest & Src
orq	Src, Dest	Dest=Dest | Src

- ```
  movq a, %rax
  movq b, %rbx
  addq %rbx, %rax
  imulq %rbx
  movq %rax, c
  ```

IMULQ has a variant with one operand which multiplies by whatever is in %rax and stores result in %rax

imulq has three forms

- imulq X : rax = X * rax
- imulq X Y : Y = X * Y
- imulq X Y Z : Z = X * Y
  ```

Some common operations with one-operand

• incq Dest	Dest = Dest + 1
• decq Dest	Dest = Dest - 1
• negq Dest	Dest = -Dest
• notq Dest	Dest = ~Dest

More Anatomy of Assembly Programs

Assembly output of hello.c

- Lines that start with "." are compiler directives.
  - This tells the assembler something about the program
  - .text is where the actual code starts.
- Lines that end with ":" are labels
  - Useful for control flow
  - Lines that start with . and end with : are usually temporary locals generated by the compiler.
- Reminder that lines that start with % are registers
- (.cfi stands for call frame information)

```
          .file    "hello.c"
          .text
          .globl   main
          .align   16, 0x90
          .type    main,@function
main:                                          # @main
          .cfi_startproc
# BB#0:
          pushq    %rbp
.Ltmp2:
          .cfi_def_cfa_offset 16
.Ltmp3:
          .cfi_offset %rbp, -16
          movq     %rsp, %rbp
.Ltmp4:
          .cfi_def_cfa_register %rbp
          subq     $16, %rsp
          leaq     .L.str, %rdi
          movl     $0, -4(%rbp)
          callq    puts
          movl     $0, %ecx
          movl     %eax, -8(%rbp)              # 4-byte Spill
          movl     %ecx, %eax
          addq     $16, %rsp
          popq     %rbp
          ret
.Ltmp5:
          .size    main, .Ltmp5-main
          .cfi_endproc

          .type    .L.str,@object              # @.str
          .section         .rodata.str1.1,"aMS",@progbits,1
.L.str:
          .asciz   "Hello Computer Systems Fall 2022"
          .size    .L.str, 33


          .ident   "clang version 3.4.2 (tags/RELEASE_34/dot2-final)"
          .section         ".note.GNU-stack","",@progbits
```

Where to Learn more?

- https://diveintosystems.org/

- Intel® 64 and IA-32 Architectures Software Developer Manuals

Document	Description
Intel® 64 and IA-32 architectures software developer's manual combined volumes: 1, 2A, 2B, 2C, 2D, 3A, 3B, 3C, 3D, and 4	This document contains the following: **Volume 1:** Describes the architecture and programming environment of processors supporting IA-32 and Intel® 64 architectures. **Volume 2:** Includes the full instruction set reference, A-Z. Describes the format of the instruction and provides reference pages for instructions. **Volume 3:** Includes the full system programming guide, parts 1, 2, 3, and 4. Describes the operating-system support environment of Intel® 64 and IA-32 architectures, including: memory management, protection, task management, interrupt and exception handling, multi-processor support, thermal and power management features, debugging, performance monitoring, system management mode, virtual machine extensions (VMX) instructions, Intel® Virtualization Technology (Intel® VT), and Intel® Software Guard Extensions (Intel® SGX). **Volume 4:** Describes the model-specific registers of processors supporting IA-32 and Intel® 64 architectures.

(Volume 2 Instruction set reference)

Bookmarks

INC—Increment by 1

Opcode	Instruction	Op/ En	64-Bit Mode	Compat/ Leg Mode	Description
FE /0	INC *r/m8*	M	Valid	Valid	Increment *r/m* byte by 1.
REX + FE /0	INC *r/m8**	M	Valid	N.E.	Increment *r/m* byte by 1.
FF /0	INC *r/m16*	M	Valid	Valid	Increment *r/m* word by 1.
FF /0	INC *r/m32*	M	Valid	Valid	Increment *r/m* doubleword by 1.
REX.W + FF /0	INC *r/m64*	M	Valid	N.E.	Increment *r/m* quadword by 1.
40+ *rw***	INC *r16*	O	N.E.	Valid	Increment word register by 1.
40+ *rd*	INC *r32*	O	N.E.	Valid	Increment doubleword register by 1.

NOTES:

* In 64-bit mode, r/m8 can not be encoded to access the following byte registers if a REX prefix is used: AH, BH, CH, DH.

** 40H through 47H are REX prefixes in 64-bit mode.

Instruction Operand Encoding

Op/En	Operand 1	Operand 2	Operand 3	Operand 4
M	ModRM:r/m (r, w)	NA	NA	NA
O	opcode + rd (r, w)	NA	NA	NA

Description

Adds 1 to the destination operand, while preserving the state of the CF flag. The destination operand can be a

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So far we looked at moving data and doing some operations on data

What's missing?

Comparisons

Compare operands: cmp_, jmp_, set__

- Often we want to compare the values of two registers
  - Think if, then, else constructs or loop exit or switch conditions

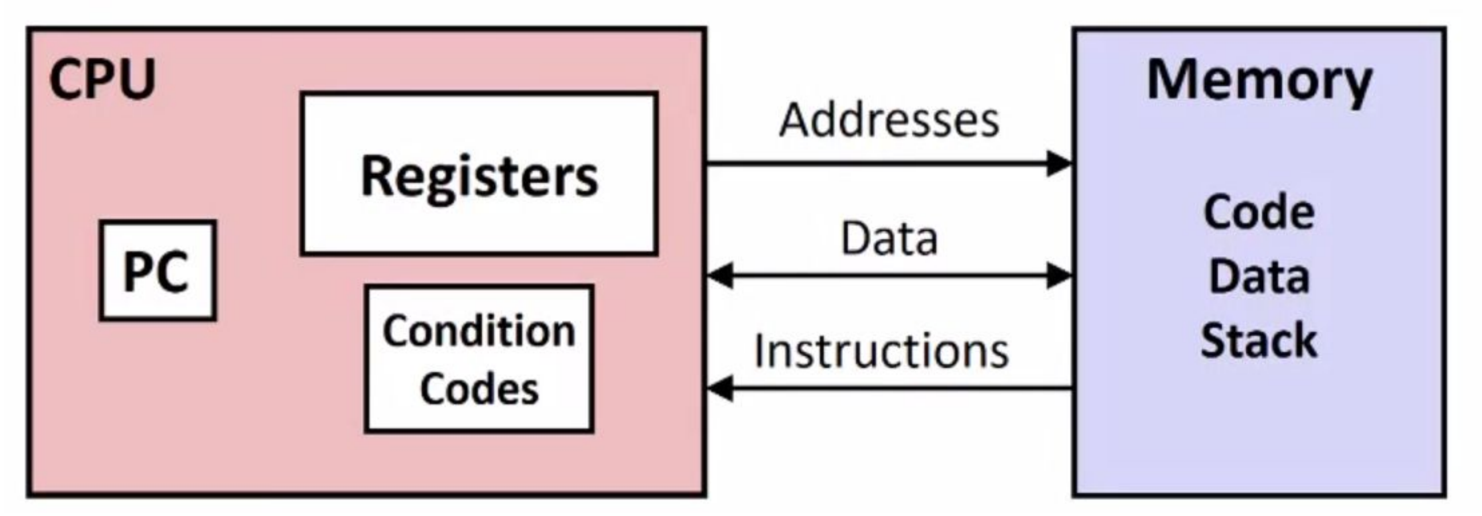
- cmpq Src2, Src1
  - cmpq Src2, Src1 is equivalent to computing Src1-Src2 (but there is no destination register)

- Now we need a method to use the result of compare, but there is not destination to find the result.

What do we do?

Remember condition codes?

- Register - where we store data (heavily used data)
- PC - gives us address of next instruction
- **Condition codes - some status information**
- Memory – where the program (code) resides and data is stored



FLAGS registers

- CF (carry flag)
  - Set to 1 when there is a carry out in an unsigned arithmetic operation
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  - Set to 1 when the result of an arithmetic operation is zero
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  - Set to 1 when there is a carry out in a signed arithmetic operation
  - Otherwise set to 0
- OF (overflow flag)
  - Set to 1 when signed arithmetic operations overflow
  - Otherwise set to 0

CS 3650 Computer Systems – Summer 1 2025

Assembly (cont.)

Unit 2

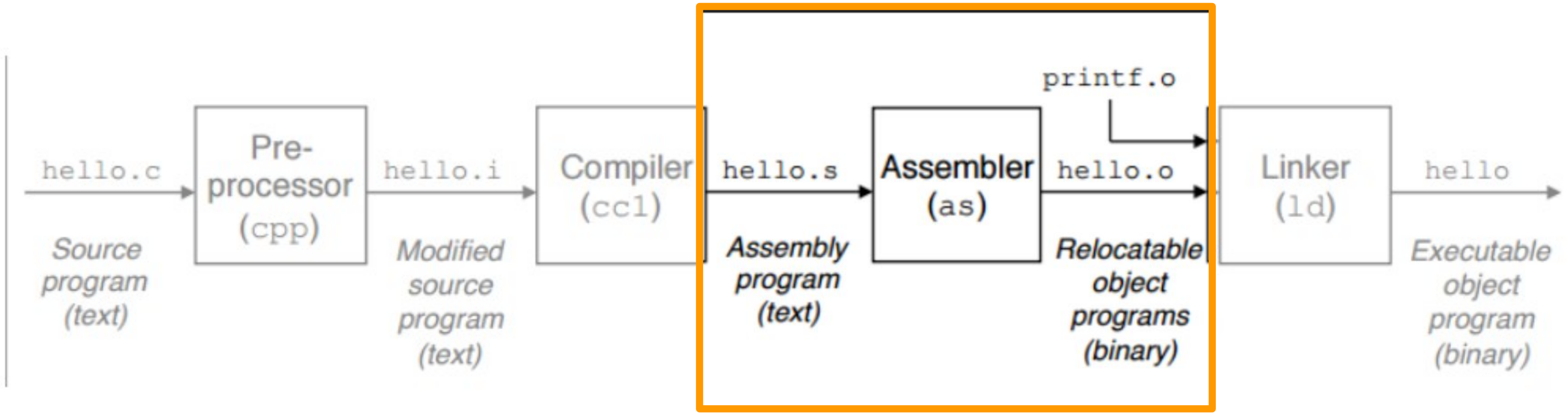
Northeastern University

* Acknowledgements: created based on Christo Wilson, Ferdinand Vesely, and Alden Jackson's lecture slides for the same course. Inspired heavily by slides from Ji-Yong Shin

Recap

Assembly is important in our toolchain

- Even if the step is often hidden from us!



Focus on this step today

- ~~Compile a program to an executable~~
  - ~~gcc main.c -o program~~

- Compile a program to assembly
  - gcc main.c -S -o main.s

- ~~Compile a program to an object file (.o file)~~
  - ~~gcc -c main.c~~

- Linker (A program called ld) then takes all of your object files and makes a binary executable.

Sizes of data types (C to assembly)

C Declaration	Intel Data Type	Assembly-code suffix	Size (bytes)
char	Byte	b	1
short	Word	w	2
int	Double word	l	4
long	Quad word	q	8
char *	Quad word	q	8
float	Single precision	s	4
double	Double Precision	l	8

x86-64 Registers

- Focus on the 64-bit column.
- These are 16 general purpose registers for storing bytes
  - (Note sometimes we do not always have access to all 16 registers)
- Registers are similar to variables where we store values

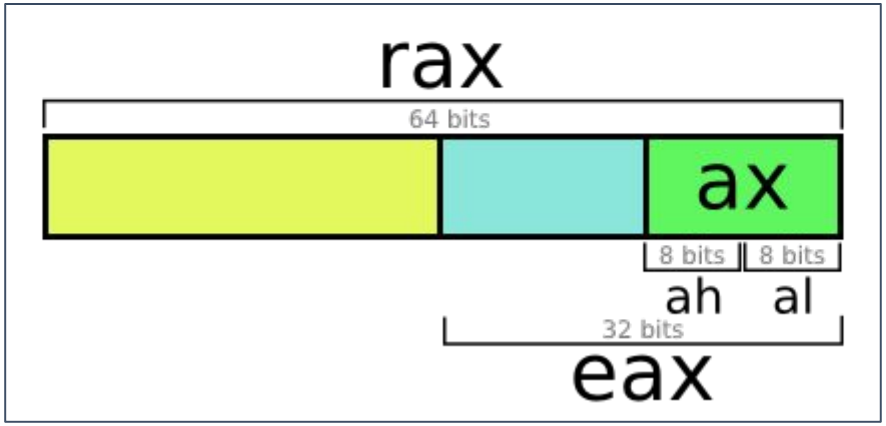
Not modified for 8-bit operands
Not modified for 16-bit operands
Zero-extended for 32-bit operands

Register encoding	Bits 63–32	Bits 31–16	Bits 15–8	Low 8-bit (bits 7–0)	16-bit	32-bit	64-bit
0			AH†	AL	AX	EAX	RAX
3			BH†	BL	BX	EBX	RBX
1			CH†	CL	CX	ECX	RCX
2			DH†	DL	DX	EDX	RDX
6				SIL‡	SI	ESI	RSI
7				DIL‡	DI	EDI	RDI
5				BPL‡	BP	EBP	RBP
4				SPL‡	SP	ESP	RSP
8				R8B	R8W	R8D	R8
9				R9B	R9W	R9D	R9
10				R10B	R10W	R10D	R10
11				R11B	R11W	R11D	R11
12				R12B	R12W	R12D	R12
13				R13B	R13W	R13D	R13
14				R14B	R14W	R14D	R14
15				R15B	R15W	R15D	R15

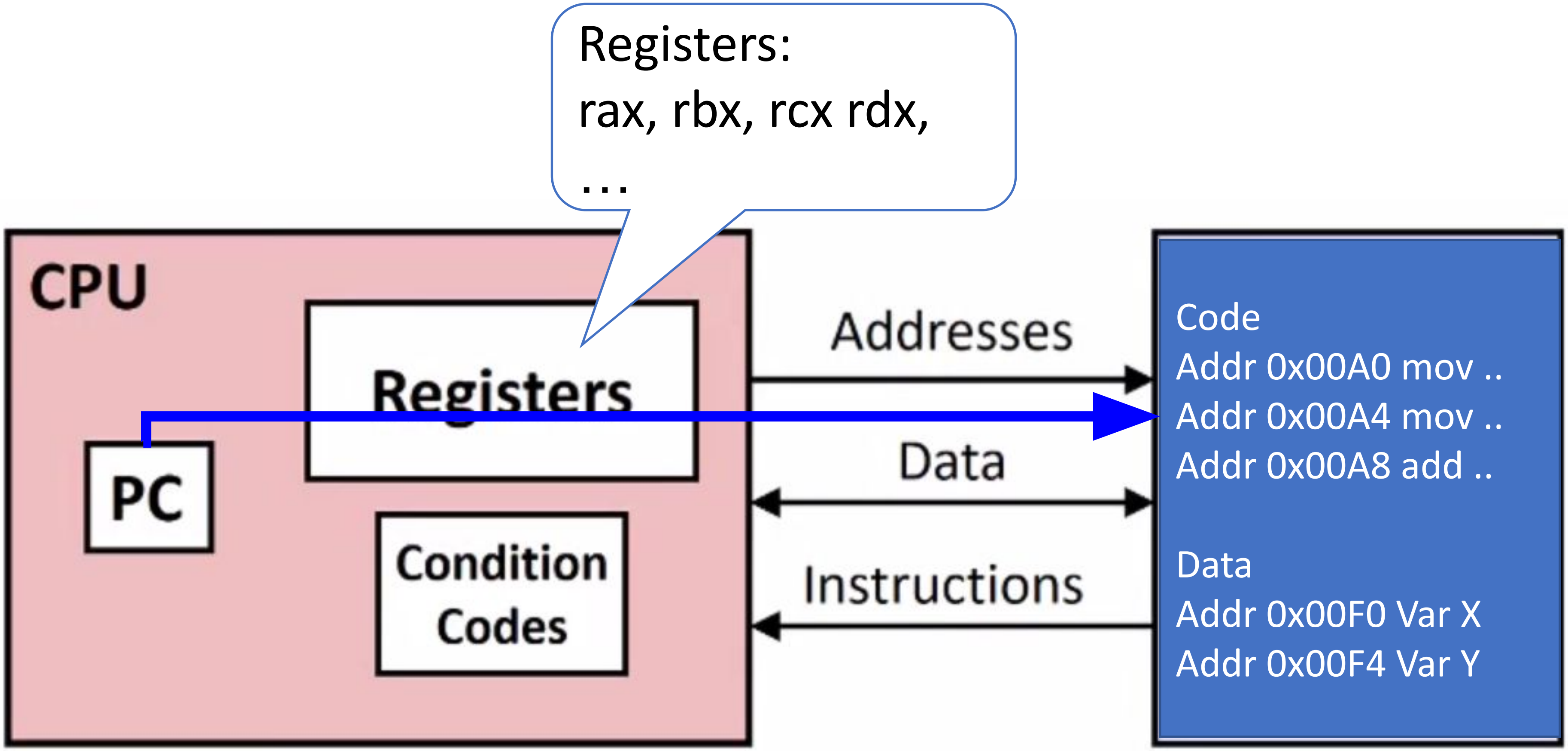
† Not legal with REX prefix ‡ Requires REX prefix

x86-64 Register (zooming in)

- Note register eax addresses the lower 32 bits of rax
- Note register ax addresses the lower 16 bits of eax
- Note register ah addresses the high 8 bits of ax
- Note register al (lowercase L) addresses the low 8 bits of ax



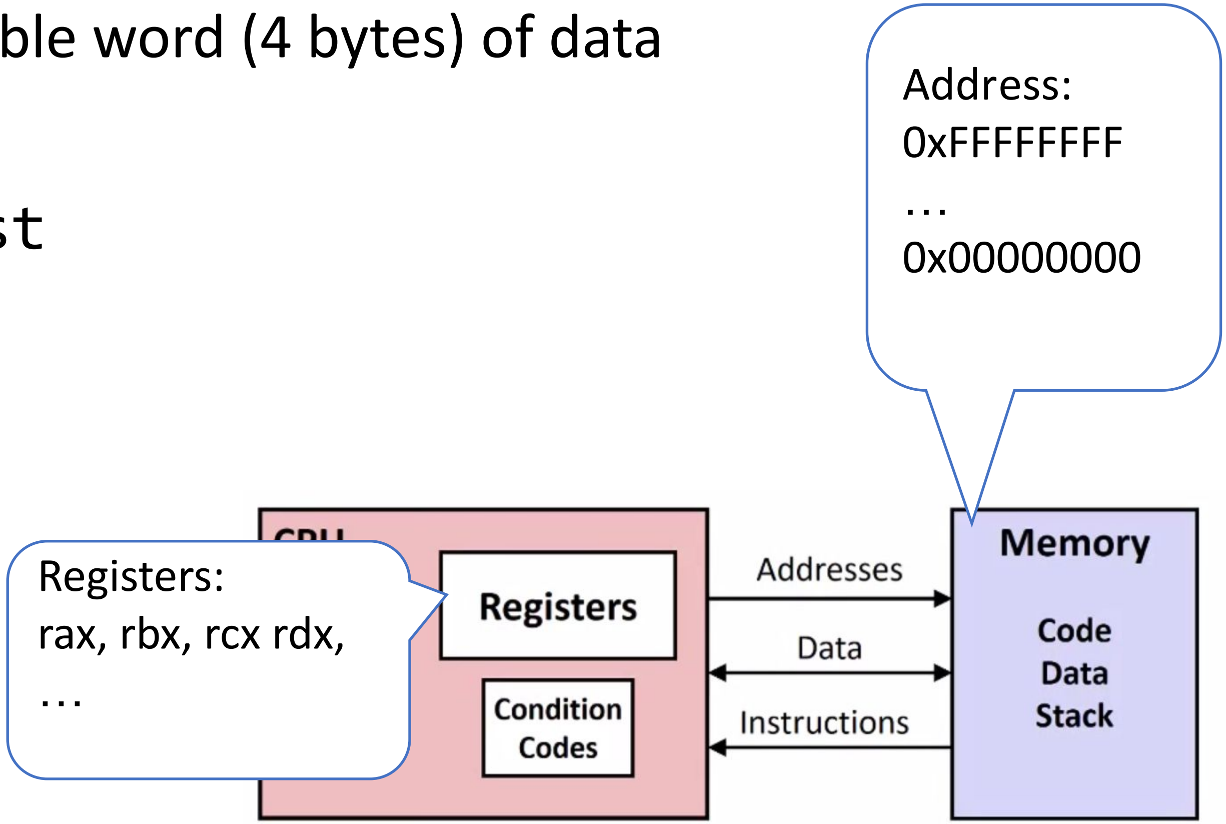
Program Counter and Memory Addresses



Moving data around | mov instruction

- (Remember moving data is all machines do!)
- movq - moves a quad word (8 bytes) of data
- movd - move a double word (4 bytes) of data

```
movq Source, Dest
```



C equivalent of movq instructions | movq src, dest

`movq $0x4, %rax`	**`%rax`** `= 0x4;` (Moving in literal value into register)
`movq $-150, (%rax)`	use value of rax as memory location and set that location to -150 (`*p = -150`)
`movq %rax, %rdx`	`%rdx = %rax` (copy src into dest)
`movq %rax, (%rdx)`	use value of rdx as memory location and set that location to value stored in rax (`*p = %rax`)
`movq (%rax), %rdx`	Set value of rdx to value of rax as memory location (`%rdx = *p`)

More assembly instructions

```
addq  Src, Dest     Dest=Dest+Src
subq  Src, Dest     Dest=Dest-Src
imulq Src, Dest     Dest=Dest*Src
salq  Src, Dest     Dest=Dest << Src
sarq  Src, Dest     Dest=Dest >> Src
shlq  Src, Dest     Dest=Dest << Src
shrq  Src, Dest     Dest=Dest >> Src
xorq  Src, Dest     Dest=Dest ^ Src
andq  Src, Dest     Dest=Dest & Src
orq   Src, Dest     Dest=Dest | Src
```

- Note on order:
  We use AT&T syntax: op Src, Dest
  Intel syntax: op Dest, Src

	Value 1	Value 2
x	0110 0011	1001 0101
x>>4 (arithmetic)	**0000** 0110	**1111** 1001
x>>4 (logical)	**0000** 0110	**0000** 1001

Compare operands: cmp_, jmp_, set__

- Often we want to compare the values of two registers
  - Think if, then, else constructs or loop exit or switch conditions

- cmpq Src2, Src1
  - cmpq Src2, Src1 is equivalent to computing Src1-Src2 (but there is no destination register)

- Now we need a method to use the result of compare, but there is not destination to find the result.

What do we do?

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  - Otherwise set to 0
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  - Set to 1 when signed arithmetic operations overflow
  - Otherwise set to 0

Conditional Branches (jumps)

Using the result from cmp => jmp instructions

- In order to read result from cmp, we use jmp to a label

Instruction		Description
`jmp`	*Label*	Jump to label
`jmp`	**Operand*	Jump to specified location
`je` / `jz`	*Label*	Jump if equal/zero
`jne` / `jnz`	*Label*	Jump if not equal/nonzero
`js`	*Label*	Jump if negative
`jns`	*Label*	Jump if nonnegative
`jg` / `jnle`	*Label*	Jump if greater (signed)
`jge` / `jnl`	*Label*	Jump if greater or equal (signed)
`jl` / `jnge`	*Label*	Jump if less (signed)
`jle` / `jng`	*Label*	Jump if less or equal
`ja` / `jnbe`	*Label*	Jump if above (unsigned)
`jae` / `jnb`	*Label*	Jump if above or equal (unsigned)
`jb` / `jnae`	*Label*	Jump if below (unsigned)
`jbe` / `jna`	*Label*	Jump if below or equal (unsigned)

Jumping to labels

```
0x8048411 <+6>:   mov    0x8(%ebp),%eax
0x8048414 <+9>:   cmp    0xc(%ebp),%eax
0x8048417 <+12>:  jle    0x8048421
0x8048419 <+14>:  mov    0xc(%ebp),%eax
0x804841f <+20>:  jmp    0x8048427
0x8048421 <+22>:  mov    0x8(%ebp),%eax
0x8048427 <+28>:  ret
```

```c
int getSmallest(int x, int y) {
    int smallest;
    if ( x > y ) {
        smallest = y;
    }
    else {
        smallest = x;
    }
    return smallest;
}
```

Jump instructions | Typically used after a compare

	Condition	Description
jmp	1	unconditional
je	ZF	jump if equal to 0
jne	~ZF	jump if not equal to 0
js	SF	Negative
jns	~SF	non-negative
jg	~(SF^OF) & ~ZF	Greater (Signed)
jge	~(SF^OF)	Greater or Equal
jl	(SF^OF)	Less (Signed)
jle	(SF ^ OF) | ZF	Less or Equal
ja	~CF & ~ZF	Above (unsigned)
jb	CF	Below (unsigned)

Conditional Branch | if-else

- long absoluteDifference (long x, long y) {

```
    long result;

    if (x > y)
        result = x-y;

    else
        result = y-x;
}
```

Take a moment to think about the ASM code

- absoluteDifference:

```
    cmpq    %rsi, %rdi
    jle     .else
    movq    %rdi,   %rax
    subq    %rsi,   %rax
    ret
.else:
    movq    %rsi,   %rax
    subq    %rdi,   %rax
    ret
```

Some reminders:

%rdi = argument x (first argument)
%rsi = argument y (second argument)
%rax = return value
cmpq src2, src1 = src1 – src2 and sets flags
jle x = jump to x if less than or equal

Code Exercise
(Take a moment to think what this assembly does)

```
    movq  $0, %rax
mystery:
    incq  %rax
    cmpq  $5, %rax
    jl    mystery
```

Code Exercise | Annotated (while loop example)

```
    movq  $0, %rax
mystery:
    incq  %rax
    cmpq  $5, %rax
    jl    mystery
```

- Move the value 0 into %rax (temp = 0)
- Increment %rax (temp = temp + 1;)
- Compare %rax with 5
- If %rax is smaller than 5 then
  jump to 'mystery'
  If not then
  proceed

Code Exercise | Annotated (while loop example)

```
    movq  $0, %rax
mystery:
    incq  %rax
    cmpq  $5, %rax
    jl    mystery
```

Equivalent C Code

```
long temp = 0;
do {
    temp = temp + 1;
  }
while(temp < 5);
```

- Move the value 0 into %rax (temp = 0)
- Label of a location
- Increment %rax (temp = temp + 1;)
- Compare %rax with 5
- If %rax is smaller than 5 then
  jump to ‘mystery’
  If not then
  proceed

Calling functions

- (Writing functions next week)
- Use call instruction
- Call accepts one operand
  - Address of function body
  - Symbolic name often used

Example:

```
call printf
```

Calling functions

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Example:

```
call printf
```

Where do arguments go?
Return values?

Calling conventions: SysV ABI x86_64

Arguments

Return value: %rax

Argument	Register
1	%rdi
2	%rsi
3	%rdx
4	%rcx
5	%r8
6	%r9

- What if there are more than six arguments?
  - Call stack

Calling printf, scanf, etc.

- Takes a variable number of arguments
- For our assignments:
  - Set %al to zero
  - `mov $0, %al`

Visit Canvas > Assignments

Work on Lab 2

Work inside login.khoury.northeastern.edu